

Attention! Because this is an open book test, the emphasis will be given to ability to understand and implement the concepts. You are allowed to use a calculator, including a graphing calculator, however the numbers were kept simple to allow calculations by hand. In all problems, please, show all important steps in you solution/proof.

1. a) (2pt) How many significant digits are in the number 03.5000?

The two definitions of accuracy are given:

Def. A. An approximation \hat{x} to x has p correct significant digits if \hat{x} and x round to the same number to p significant digits. Rounding is the act of replacing a given number by the nearest p significant digits number with some rule for breaking ties when the are two nearest.

Def. B. \hat{x} agrees with x to p significant digits if $|x - \hat{x}|$ is less than half a unit in the pth significant digit of x.

b) (4pt) To how many significant digits does 3.4991 approximate 3.5000 according to the Def. A

c) (4pt) To how many significant digits does 3.4991 approximate 3.5000 according to the **Def. B**

2. a) (10pt) Perform Rounding Error analysis for the problem of computing determinant of a 2×2 matrix in floating point arithmetic:

$$d = \det \left(\begin{array}{cc} a_{11} & a_{12} \\ a_{21} & a_{22} \end{array} \right).$$

Assume that the calculations in the determinant formula $\hat{d} = fl(\det(A)) := fl(a_{11}a_{22} - a_{12}a_{21})$ are done from left to right.

b) (5pt) Based on your answer in part a), is it possible to derive a backward error result for computing the quantity d = det(A) in the form

$$\hat{d} = \det(A + \Delta A)?$$

3. (15pt) Prove that the subordinate matrix norm for the vector norm $||x||_{\infty} = \max_{1 \le i \le n} |x_i|$ is given by the formula $(A \in \mathbb{R}^{m \times n})$

$$||A||_{\infty} = \max_{1 \le i \le m} \sum_{j=1}^{n} |a_{ij}|.$$

4. (15pt) (Refers to page 123) Use Theorem 7.4 to prove formula (7.11).

Hint. Recall Theorem 7.4:

Theorem 7.4. Let Ax = b and $(A + \Delta A)y = (b + \Delta b)$, where $|\Delta A| \leq \epsilon E$ and $|\Delta b| \leq \epsilon f$, and assume that $\epsilon ||A^{-1}|E|| < 1$, where $||\cdot||$ is an absolute norm. Then

$$\frac{\|x-y\|}{\|x\|} \le \frac{\epsilon}{1-\epsilon\| \, |A^{-1}|E\,\|} \frac{\| \, |A^{-1}|(E|x|+f)\|}{\|x\|},$$

and for the ∞ -norm this bound is attainable to first order in ϵ .

Recall the definition of the condition number:

$$\operatorname{cond}_{E,f}(A,x) := \lim_{\epsilon \to 0} \sup \left\{ \frac{\|\Delta x\|_{\infty}}{\epsilon \|x\|_{\infty}} : (A + \Delta A)(x + \Delta x) = b + \Delta b, |\Delta A| \le \epsilon E, |\Delta b| \le \epsilon f \right\}$$

Use Theorem 7.4 to prove that

$$\operatorname{cond}_{E,f}(A, x) = \frac{\| |A^{-1}|(E|x| + f)\|_{\infty}}{\|x\|_{\infty}}$$

Final. Name:_____ ID: _____

5. a) (7pt) (Refers to page 142) Apply Lemma 8.4 to the problem of solving an upper triangular system Ux = b by substitution: $x_i = (b_i - \sum_{j=i+1}^n u_{ij}x_j)/u_{ii}$.

Hint. Recall Lemma 8.4.: Lemma 8.4. If $y = (c - \sum_{i=1}^{k-1} a_i b_i)/b_k$ is evaluated in floating point arithmetic, then, no matter what the order of evaluation,

$$b_k \hat{y}(1+\theta_k^{(0)}) = c - \sum_{i=1}^{k-1} a_i b_i (1+\theta_k^{(i)}),$$

where $|\theta_k^{(i)}| \leq \gamma_k$ for all *i*. If $b_k = 1$, so that there is no division, then $|\theta_k^{(i)}| \leq \gamma_{k-1}$ for all *i*.

Recall the substitution formula:

$$x_i = (b_i - \sum_{j=i+1}^n u_{ij} x_j) / u_{ii}$$

Apply Lemma 8.4 to the substitution formula

b) (8pt) Use the results of part a) to prove Theorem 8.5 in the case when T is upper triangular:

Theorem 8.5 Let the triangular system Tx = b, where $T \in \mathbb{R}^{n \times n}$ is nonsingular, be solved by substitution, with any ordering. Then the computed solution \hat{x} satisfies

$$(T + \Delta T)\hat{x} = b, \qquad |\Delta T| \le \gamma_n |T|.$$

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6. For the matrix

$$A = \left(\begin{array}{rrr} 1 & 1 & 0\\ 0 & \varepsilon & \varepsilon\\ 0 & 0 & 1 \end{array}\right)$$

Calculate the following in ∞ -norm: a) (5pt) $\kappa(A)$. Hint:

$$A^{-1} = \begin{pmatrix} 1 & -1/\varepsilon & 1\\ 0 & 1/\varepsilon & -1\\ 0 & 0 & 1 \end{pmatrix}$$

b) (5pt) $\operatorname{cond}(A)$.

7. (10pt) On the 3rd stage of Gaussian elimination with Complete Pivoting the following matrix has been obtained:

$$A^{(k)} = \begin{pmatrix} -7 & 0.3 & 0.1 & -0.2 & 0.3 & 4 \\ 0 & 0.5 & 0.1 & 0.3 & 0.1 & -0.2 \\ 0 & 0 & 2 & 0.2 & 0.6 & 0.5 \\ 0 & 0 & 0.5 & 3 & -5 & -3 \\ 0 & 0 & 0.4 & -0.3 & -2 & 1 \\ 0 & 0 & 1 & -0.1 & 4 & 0.1 \end{pmatrix}$$

What rows are and what columns have to be permuted to calculate the next step.

8. (15pt) (Refers to page 198) Use Theorem 10.3. to prove Theorem 10.4.:

Theorem 10.4. Let $A \in \mathbb{R}^{n \times n}$ be symmetric positive definite and suppose Cholesky factorization produces a computed factor \hat{R} and a computer solution \hat{x} to Ax = b. Then,

$$(A + \Delta A)\hat{x} = b, \qquad |\Delta A| \le \gamma_{3n+1} |\hat{R}^T| |\hat{R}|.$$

Hint. Use Theorem 10.3.:

Theorem 10.3. If Cholesky factorization applied to the symmetric positive definite matrix $A \in \mathbb{R}^{n \times n}$ runs to completion then the computed factor \hat{R} satisfies

$$|\hat{R}^T||\hat{R}| = A + \Delta A, \qquad |\Delta A| \le \gamma_{n+1}|\hat{R}^T||\hat{R}|.$$

Recall that if $A = R^T R$, then the linear problem Ax = b is equivalent to $R^T Rx = b$ which can be solved in two steps: $R^T y = b$ and Rx = y.

Use proof of Theorem 9.4 as the example.